How to design a DRAM Controller to interface a DRAM with the SHARC DSP

**Introduction:**
This document provides a reference design for customers on how to interface a DRAM with the SHARC DSP using a DRAM Controller. The controller is implemented as a state-machine on a Xilinx PLD. This design has been simulated, built in hardware, and tested to operate successfully. The design, however, has not been optimized for best performance. Recommendations for further modifications as well as alternative designs will be given in Section IV.

**I. Design Overview:**
Different applications require specific ways of DRAM interfacing. This design implements a DRAM controller to interface a 4 Meg DRAM by Micron with the SHARC.

![Diagram of DRAM Interfacing](image)

**Figure 1: DRAM Interfacing Block Diagram**
DRAM is an ideal solution for mass data storage with high speed. However, DRAM requires page swapping if an access crosses a page boundary. In addition, DRAM requires refreshing so that data is retained properly. This means an external controller is needed to handle the interface between the Sharc processor and the DRAM.

The controller receives input signals from the Sharc on each memory access and asserts the right control signals to the DRAM to read or write. Eleven 2:1 mux's are used to select between column address and row address depending on whether the controller is doing a page swap or a normal read/write. The refresh timer is responsible for telling the controller when to refresh the DRAM.

II. Descriptions of the Refresh Timer and the Controller

1. Refresh Timer:

Data sheets for this 4 Meg DRAM specify that all 2048 rows or pages need to be refreshed every 32ms. Spreading this task into equal intervals shows that a new row needs to be refreshed every 15.625µs (32ms/2048). In this design, a refresh request is generated by the timer every 15.525µs or 621 cycles if using a 40 Mhz clock.

The refresh timer works as follows: during reset, the counter is forced to zero from any state. When reset is released, the counter will start counting up to 620. At this point, the counter resets to zero and starts counting again. At the same time, a refresh request (RFQ) is asserted, signaling the controller to refresh one row of the DRAM. There are different ways of refreshing the DRAM. For this design, CBR or CAS Before RAS method is used as illustrated in the diagram below. First, the controller returns CAS and RAS to high, if necessary. Next, the controller brings CAS low, and then RAS low in the next cycle. It is important that RAS and CAS should remain high or low long enough to meet the time specifications in the data sheets. At the end of the refresh cycles, the controller asserts RFC (refresh complete) to signal the refresh timer to bring RFQ low.

Using the CBR refresh method, the user does not have to worry about which row to refresh. The DRAM itself has an internal counter that keeps track of which row to refresh. Each time a CBR is performed, this internal counter will increment by one so that the next adjacent page is refreshed on the next CBR.

By refreshing one row every 15.5µs, the entire 4 Meg DRAM will effectively be refreshed within every 32ms.

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Figure 2: Dram Refreshing for One Row or Page
2. DRAM Controller:

a. Normal Page Swap, Read, Write

Figure 3 below illustrates the operations of an out-of-page read followed immediately by a within-page write. Note that the timing diagram given here does not reflect precisely the number of cycles used in the Abel file. Refer to discussion about optimization toward the end of this document.

Upon the detection of /MS and Page, the controller will assert MuxSel to select the row address. Then /RAS is asserted low to latch in the new row address. Note that MuxSel is clocked at the negative clock edge so that row address is valid by the falling edge of /RAS. Then, if this is a read, /WR_DRAM is kept high while /CAS is asserted low. During the same cycle, Ack is asserted high to indicate the completion of one memory access. Note that /RAS is held low so that the same page is retained. It will remain low until a page swap is required or a DRAM refresh is serviced.

In this example, a write follows immediately. Since it is a same-page operation, only /CAS needs to go low. /WR_DRAM is asserted this time for a write cycle.

b. Refresh Request during a Read/Write:

If the DRAM controller is servicing a memory access when a refresh request occurs, the controller will keep going to complete that memory read or write before entering the refresh cycles. RFQ will stay high until one row refresh has been done and RFC is asserted. This is illustrated in Figure 4.

c. Memory Access Request during a Refresh Service:

If the Sharc asserts /MS while the controller is refreshing the DRAM, the controller will continue to finish the refresh before servicing the memory access. Without having Ack high during this time, the Sharc will keep driving /MS, /RD or /WR, Data and Address lines. Note that in this design, the Sharc is configured to use Ack-only as waitstates. Figure 5 illustrates this scenario.

![Figure 3: Page Swap, Read, Write](Image)
It's time to refresh!

**Figure 4: Refresh Request during a Read/Write**

It's time to refresh!

**Figure 5: Memory Access Request during a Refresh Service**
d. Controller State Machine:

This DRAM controller is implemented as a Mealy state machine and is programmed onto a Xilinx PLD with 69 I/O pins. The state diagram on the following page shows the flow of the design with transitions from one state to another. Please refer to the Abel source code for the output signals.

Several states in the state machine are used merely to hold /CAS or /RAS low or high in order to meet the time specifications. Further optimization could be done to speed up the design as discussed in a later section.

During the Idle/Reset cycle, the controller keeps /RAS and /CAS high which will put the DRAM into a low power mode. From here, the controller can go to refresh cycles if RFQ is asserted or go to start a memory access. Note that coming out of the idle state, the controller needs to update the page in the DRAM regardless of whether this is a out-of-page access or not. The reason is that /RAS is no longer held low during idle period.

In the state of Read/Write within Page, the controller can remain here to wait for the next memory access or go to refresh the DRAM if necessary.

**Figure 6: DRAM Controller State Diagram**
III. Hardware Description
A prototype of this design has been built with the layout shown in Figure 7 below.

On the left is the DRAM board with 2 major components:

1. DRAM: This is placed along the long edge close to the Sharc EZ-Kit Lite board so that the lengths of the wires connected to the Sharc EZ-Kit Lite board are minimized for noise reduction purposes.
2. Xilinx PLD: This is the DRAM Controller. The refresh timer and the muxes are all programmed onto this same PLD. The Abel file attached with this document includes all code and equations necessary to program this PLD.

Other components include the following:

1. A JTAG connector that allows the users to easily program the PLD using a parallel EZJTAG download cable from Xilinx.
2. An optional socket to hold the Xtal if an external clock is used.
3. Power connector. An external power source that can provide at least 1.5A must be used to satisfy the demand of the DRAM.

On the right is the Sharc EZ-KIT Lite board with four external connectors shown in the figures. All the signals shown in Figure 1: DRAM Interfacing Block Diagram are available from these connectors. This board runs on its own power supply. However, common ground is needed between the two boards. At least one thick wire should be used to make this connection between the ground plane of the DRAM board and the ground pin of the power connector to the Sharc EZ-KIT Lite board.

Twisted pairs of wires are used for all data lines, address lines, and control signals with the looping wire grounded to reduce signal interference. Same method is used for the clock signal but with thicker wire.

Figure 7: Components and Board Layout

Figure 8: Twisted Pair of Wires
Each VDD pin on the DRAM and the PLD has two decoupling capacitors of .02uF and 1.0uF in parallel. On the pin where the power supply is brought into the board, two capacitors of .1uF and 22uF are used.

Note that on the Sharc EZ-KIT Lite board, there is an EZ-ICE connector which allows the users to attach an EZ-ICE and use this emulator to debug the DRAM Controller.

IV. Possible Design Modifications and Improvement

*** Important Notes about This Design ***

However, the DRAM Controller has only been tested with a 40Mhz clock taken from the Sharc EZ-KIT Lite board. This means there are two possible improvements for speed.

1. **60 Mhz clock is used:**
   In this Design we use an external EPSON Americana 60 Mhz Oscillator(SG-615PCV) in addition to the other hardware to make the controller function faster. In the Appendix of this EE Note we provide
   
   App. 1. **Abel Source Code:** consists of equations to implement the refresh timer, mux, and the state machine for the controller. The Abel code has been synthesized and fitted using the Foundations Series Tool Suite
   
   App. 2. **Schematics:** is drawn using OrCAD. J1, J3, and J6 are the actual connectors on the Sharc EZ-Kit Lite board.
   
   App. 3. **Layout**
   
   App. 4. **Overall Layout**

2. **66Mhz or higher clock is used:**
   If a faster external clock is used as originally intended for this design, then other considerations have to be taken into account. The immediate benefit would be speed. However, there will be timing issues such as making Ack meet the setup and hold time so that the Sharc can sample correctly. Another potential problem is that running a 66Mhz will introduce more noise to the system, especially for a wire-wrapping board.

   It is also recommended to optimize the logic of the state-machine for better performance, especially those states that extend the /CAS and /RAS control signals.

   Of course, another solution to this design problem would be changing the entire approach to a more effective method. If the controller is run at a fast clock speed, it is possible just to program waitstates on the SHARC and not use the ACK signal at all. If a refresh request occurs during a burst of write or read cycles, the controller may assert the Suspend Bus Tri-State (SBSTS) signal to halt the dsp, and then service the refresh. This way the SHARC will not occupy the address bus and data bus while the controller is refreshing the DRAM. If there are other external devices to the system, these devices may make use of the bus while the dsp is halting.

   /SBSTS may also be used during a page swapping cycle. The controller latches in the new row address before asserting /SBTS. When the address lines of the dsp become tri-stated, the controller puts back the row address onto the bus and asserts /RAS to update the page in the DRAM. Then, it releases /SBTS to continue the memory access normally.

V. Conclusion

This design has been tested to work quite successfully. However, this is not the optimal design as far as speed is concerned. Once in a while, it is still subjective to noise, especially when all or most of the data lines switch to high at the same time causing a few glitches on some control signals. If someone tries to duplicate this process, it is recommended that further measures be taken to clean out such noises. For example, putting a capacitor and a resistor (with calculated values) in series near input pins to the Sharc might help reduce reflection on these transmission lines.
MODULE DramCtrl

title 'Dram Controller';

"Input signals"

Reset_   pin 81   istype 'reg';
DramClkin pin 9   istype 'reg'; "Ext 60MHz Cloc
k
MS0_      pin 53   istype 'reg';
RD_       pin 36   istype 'reg';
WR_       pin 35   istype 'reg';
Page      pin 33   istype 'reg';
A00,A01,A02,A03,A04 pin 80,79,78,77,76,75 istype 'com';
A05,A06,A07,A08,A09,A10 pin 74,72,71,70,69,68 istype 'com';
A11,A12,A13,A14,A15 pin 67,66,65,63,62 istype 'com';
A16,A17,A18,A19,A20,A21 pin 61,58,57,56,55,54 istype 'com';

"Input/Output"

RFQ      pin 50   istype 'reg'; "Timer's Outp
ut to Controller's Input"
RFC      pin 51   istype 'reg'; "Controller's
Output to Timer's Input"

"Output signals"

RAS_      pin 24   istype 'reg';
CAS_      pin 25   istype 'reg';
WR_DRAM_  pin 26   istype 'reg';
Ack       pin 34   istype 'reg'; "Ack needs to
be in sync with SHARC Clk
MuxSel    pin 52   istype 'reg'; "H: RowAddr; L
: ColAddr

Y00,Y01,Y02,Y03,Y04 pin 11,12,13,14,15 istype 'com';
Y05,Y06,Y07,Y08,Y09,Y10 pin 17,18,20,21,23,19 istype 'com';
q0, q1, q2, q3 pin 43,44,45,46 istype 'reg'; "Variables for
State Machine
t0, t1, t2, t3, t4 pin 82,83,84,1,2 istype 'reg';
t5, t6, t7, t8, t9 pin 3,4,5,6,7 istype 'reg';

xepld property 'FAST ON'; "Special Setting for Xilinx CPLD

declarations

S0 = [0,0,0,0]; S1 = [0,0,0,1]; S2 = [0,0,1,0]; S3 = [0,0,1,1];
S4 = [0,1,0,0]; S5 = [0,1,0,1]; S6 = [0,1,1,0]; S7 = [0,1,1,1];
S8 = [1,0,0,0]; S9 = [1,0,0,1]; S10 = [1,0,1,0]; S11 = [1,0,1,1];
S12 = [1,1,0,0]; S13 = [1,1,0,1]; S14 = [1,1,1,0]; S15 = [1,1,1,1];

*Intermediate variables
"use the following line for actual code : cnt_ref = 930;
CNT_REF = t9 & t8 & t7 & !t6 & t5 & !t4 & !t3 & !t2 & t1 & !t0;
*use the following line for simulation : cnt_ref = 31;
*CNT_REF = !t9 & !t8 & !t7 & !t6 & !t5 & !t4 & !t3 & t2 & t1 & t0;
count = [t9,t8,t7,t6,t5,t4,t3,t2,t1,t0];
lowerAddr = [A10,A09,A08,A07,A06,A05,A04,A03,A02,A01,A00];
upperAddr = [A21,A20,A19,A18,A17,A16,A15,A14,A13,A12,A11];
fsrState = [q3,q2,q1,q0];
SCLR  = !Reset_ # CNT_REF;
C,H,L,Z  = .C.,1,0,.Z.;

equations
Dramctrl

[q3,q2,q1,q0].clk = DramClkin;
[RAS_,CAS_,WR_DRAM_,Ack].clk = DramClkin;
[RFQ,RFC].clk = DramClkin;
[MuxSel].clk = !DramClkin;
[t0,t1,t2,t3,t4,t5,t6,t7,t8,t9].clk = DramClkin;

Ack.oe=Reset_; /* This is used to disable Ack when EZ-Kit Lite is booting.
This reset is separate from Sharc Reset. This is a temporary
work around. So that the UART and the DRAM don't conflict.

*Below is the refresh timer:
t0 := (!t0) & !SCLR;
t1 := (t1 $ t0) & !SCLR;
t2 := (t2 $ (t1 & t0)) & !SCLR;
t3 := (t3 $ (t2 & t1 & t0)) & !SCLR;
t4 := (t4 $ (t3 & t2 & t1 & t0)) & !SCLR;
t5 := (t5 $ (t4 & t3 & t2 & t1 & t0)) & !SCLR;
t6 := (t6 $ (t5 & t4 & t3 & t2 & t1 & t0)) & !SCLR;
t7 := (t7 $ (t6 & t5 & t4 & t3 & t2 & t1 & t0)) & !SCLR;
t8 := (t8 $ (t7 & t6 & t5 & t4 & t3 & t2 & t1 & t0)) & !SCLR;
t9 := (t9 $ (t8 & t7 & t6 & t5 & t4 & t3 & t2 & t1 & t0)) & !SCLR;

RFQ := !RFC & ((CNT_REF # RFQ.FB) & Reset_);

*Shown here is the Multiplexer for Row and Column Addresses.
Y00 = (A00 & !MuxSel ) # (A11 & MuxSel);
Y01 = (A01 & !MuxSel ) # (A12 & MuxSel);
Y02 = (A02 & !MuxSel ) # (A13 & MuxSel);
Y03 = (A03 & !MuxSel ) # (A14 & MuxSel);
Y04 = (A04 & !MuxSel ) # (A15 & MuxSel);
Y05 = (A05 & !MuxSel ) # (A16 & MuxSel);
Y06 = (A06 & !MuxSel ) # (A17 & MuxSel);
Y07 = (A07 & !MuxSel ) # (A18 & MuxSel);
Y08 = (A08 & !MuxSel ) # (A19 & MuxSel);
Y09 = (A09 & !MuxSel ) # (A20 & MuxSel);
Y10 = (A10 & !MuxSel ) # (A21 & MuxSel);

state_diagram [q3,q2,q1,q0];

**************************Idle/Reset State**************************
State S0:
IFT(!Reset_)
THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (RFQ)
THEN S1 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (!MS0_)
THEN S8 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE
Goto S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}

**************************Refreshing states**************************
State S1:
IFT(!Reset_)
THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE
Goto S2 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
State S2:
IFT(!Reset_)
THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}

App. 1-2
ELSE
  Goto S3 WITH (RAS_ := 1; CAS_ := 0; WR_DRAM_ := 1;)

State S3:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S4 WITH (RAS_ := 0; CAS_ := 0; WR_DRAM_ := 1;)

State S4:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S5 WITH (RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1;)

State S5:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S6 WITH (RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1;)

State S6:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S7 WITH (RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1; RFC := 1;)

State S7:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE IF (!MS0_)
    THEN S8 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)

"*****************PageSwap State**************************

State S8:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S9 WITH (MuxSel := 1; RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)

State S9:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S10 WITH (MuxSel := 1; RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)

State S10:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S11 WITH (MuxSel := 1; RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1;)

State S11:
  IF (!Reset_)
    THEN S0 WITH (RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;)
  ELSE
    Goto S12 WITH (MuxSel := 1; RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1;)

"*****************Waiting for RD/WR State after PageSwap************
IF (!Reset_)
    THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (!RD_)
    THEN S14 WITH {RAS_ := 0; CAS_ := 0; WR_DRAM_ := 1; Ack := 0;}
ELSE IF (!WR_)
    THEN S14 WITH {RAS_ := 0; CAS_ := 0; WR_DRAM_ := 0; Ack := 0;}
ELSE
    Goto S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}

"*****************Waiting for RFQ or next mem access***************
State S13:
IF (!Reset_)
    THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (RFQ)
    THEN S1 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (!MS0_ & Page)
    THEN S8 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (!MS0_ & !Page & !RD_)
    THEN S14 WITH {RAS_ := 0; CAS_ :=0; WR_DRAM_ := 1; Ack := 0;}
ELSE IF (!MS0_ & !Page & !WR_)
    THEN S14 WITH {RAS_ := 0; CAS_ :=0; WR_DRAM_ := 0; Ack := 0;}
ELSE
    Goto S13 WITH {RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1;}

"******************Extending CAS/RAS State*************************
State S14:
IF (!Reset_)
    THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE IF (!RD_)
    THEN S15 WITH {RAS_ := 0; CAS_ :=1; WR_DRAM_ := 1; Ack := 1;}
ELSE IF (!WR_)
    THEN S15 WITH {RAS_ := 0; CAS_ :=1; WR_DRAM_ := 1; Ack := 1;}
ELSE
    GOTO S0 WITH {RAS_ := 0; CAS_ :=1; WR_DRAM_ := 1;}

State S15:
IF (!Reset_)
    THEN S0 WITH {RAS_ := 1; CAS_ := 1; WR_DRAM_ := 1;}
ELSE
    GOTO S13 WITH {RAS_ := 0; CAS_ := 1; WR_DRAM_ := 1; Ack := 0;}
END DramCtrl
Controller can be set in reset state while the SHARC EZKIT LITE uses the ACK signal. In other words, the ACK of UART will not conflict with Jumper is placed so that the Controller can be set with reset of EZKIT LITE. EZ-ICE program will still work; simply connect Reset of Controller with ACK of Controller. It is recommended to use the ACK output of UART to control the Jumper. In order to avoid the Ack of UART to conflict with ACK signal.
Please look at Table for exact pinout connections.

Connected to CPLD

Connected to SDRAM

Ground Plane
PLEASE LOOK AT TABLE FOR EXACT PINOUT CONNECTIONS.
SHARC EZ-Kit Lite is upside down on bottom

Top Board is the Proto Board facing up.